

# CL13: Big-O Notation

## Announcements

#### • Exercise 4 - List Utils

- "extend" should not be using in built-in extend method!
- I'll extend the deadline through *tomorrow* for you to make changes
- After you re-submit exercise, submit a regrade request!
- Quiz 03
  - Multiple Choice: Big-O Runtime, Unit Tests, Dictionaries
  - Memory Diagram: Dictionaries, Functions
  - Short Response: Unit Tests and Dictionary Syntax
  - Function Writing: Dictionaries
  - *Monday's class will be a virtual review session!!!*

## **Recall: Algorithms**

**Input** is data given to an algorithm

An algorithm is a series of steps

An algorithm **returns** some **result** 

An algorithm *may* be influenced by its **environment** and it *may* produce side-effects which influence its environment.



## What is an algorithm?

- A set of steps to solve a general problem
- Finite
- Can handle a problem of arbitrary size

## How do we measure how "good" an algorithm is?

- Is it correct?
- How long does it take to implement?
- How much computer memory does it take?

## Why do we care about computation speed?

- Security: Cryptography works because encrypted information takes *too long* to decipher!
- User Experience: Users don't want to work with a slow application!
- Big Data: We want to be able to feed as much data as possible into our systems, but we need a way to *efficiently* do that!

### Measurements We Use

- *O* "Big O": upper bound (worst case runtime)
- $\Omega$  "Big Omega": lower bound (best case runtime)
- 🕞 "Big Theta": average runtime

## Returning to Finding the Lowest Card in a Deck



- Go from left to right
- Remember the lowest card you've seen *so far* and compare it to the next cards































2













2





















4 actions for input of 4 cards.













4 actions for input of 4 cards.  $\rightarrow$ n actions for input of size n.

• In this approach, we always have to check every card in the deck, so our runtime will always be approximately *n* where *n* is the size of the deck.

Finding the minimum  $\in O(n)$ Finding the minimum  $\in \Omega(n)$ Finding the minimum  $\in \Theta(n)$ 

## Speed vs. Memory

- Sometimes you can make a tradeoff between speed and memory.
- E.g. storing a value rather than computing it repeatedly.

```
1
     def find_min1(nums: list[int]) -> int:
         min_idx: int = 0
 2
         idx: int = 0
 3
         while idx < len(nums):
 4
             if nums[idx] < nums[min_idx]:</pre>
 5
                 min_idx = idx
 6
 7
             idx += 1
 8
         return min_idx
 9
10
     def find_min2(nums: list[int]) -> int:
11
         min_idx: int = 0
12
         min_val: int = nums[min_idx]
         idx: int = 0
13
14
         while idx < len(nums):</pre>
15
             val: int = nums[idx]
             if val < min_val:</pre>
16
                 min_idx = idx
17
18
                 min val = val
19
             idx += 1
20
         return min_idx
21
22
     search_vals: list[int] = [10, 9, 8]
     find_min1(search_vals)
23
24
     find_min2(search_vals)
```

## New Example: Finding a specific card.



- Go from left to right
- The first time you see your card, exit!

# Finding 3









# Finding 3









# Finding 3











What is the worst case input for this algorithm? (What will make us look at the *most* cards before exiting?)

What is the Big-O (worst case) runtime in terms of deck size *n*?